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Parallel Ant Colony Algorithm for Shortest Path Problem

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Abstract

During travelling, more and more information must be taken into account, and travelers have to make several complex decisions. In order to support these decisions, IT solutions are unavoidable, and as the computational demand is constantly growing, the examination of state-of-the-art methodologies is necessary. In our research, a parallelized Ant Colony algorithm was investigated, and a parameter study on a real network has been made. The aim was to inspect the sensibility of the method and to demonstrate its applicability in a multi-threaded system (e.g. Cloud-based systems). Based on the research, increased effectiveness can be reached by using more threads. The novelty of the paper is the usage of the processors' parallel computing capability for routing with the Ant Colony algorithm.

Keywords

shortest path, parallel computing, multimodal transportation network, Ant Colony algorithm, super network, cloud computing

1 Introduction

The multimodal network should contain a certain number of data that the algorithms have to process in the shortest possible time. This paper aims to lay down the foundation for our further work in a large-sized multimodal network.

In our previous research paper, the weaknesses of traditional algorithms were introduced [1]. Their main drawback is that they produce only one solution per one parameter set and they do not give alternative solutions (multiple routes). Using an artificial intelligence method - in our case, Ant Colony algorithm - is a possible way to calculate multiple alternative options; moreover, this system is capable of taking multiple input parameters (criteria) into account. The drawback is: Ant Colony is a resource-intensive method, however, the parallelization of the method could be expedient on large networks.

More and more researchers are involved in this field, as time saving, better efficiency and cost optimizations are getting more important. In 2012, a paper deeply analyzing this topic was published at the Eindhoven University [2]. They introduced the multimodal routing system. In this paper, the Dijkstra [3] [4] and A^* [4] [5] algorithms were used. The conclusion was that there is some limitation in large networks. To solve this problem, it is possible to use caching techniques or to limit the search area, which was introduced in a paper where a super network was used. [6]

Dominik Kirscher used the Ile-de-France region in France in his research [7], and his aim was to find a route with the same origin and destination. This system is perfect for commuters who use P+R facilities every day, and go in the city center with public transport; and, at the end of the day, they arrive back to the same P+R place. In this research, the SDALT algorithm was used instead of Dijkstra to speed up the calculations.

The old traveling salesman problem (TSP) and its solutions shall not be forgotten either. One of the solutions is the Branch and Bound Algorithm. G. Laporte and Y. Nobert made research with this method [8]. Their aim was to apply the algorithm on capacitated vehicles with high number of cities. This development was continued in 2003 by Jens Lysgaard, Adam N. Letchford and Richard W. Eglese [9]. The original method was supplemented by cut pool management to increase its efficiency. Some years later, the robustness was increased [10].

A genetic algorithm based solution was implemented by Haicong Yu and Feng Lu [11]. Their system, with an improved genetic algorithm was capable of handling a multimodal network. They showed that the genetic algorithm has the ability to investigate the options simultaneously.

A Helsinki commuter's network was investigated by 3 scientists [12]. An application was developed where the public transport and P+R facilities were linked together. As a result, the level of CO_2 emission was decreased and P+R usage was increased.

A Hungarian survey was made to compare the available route planners in Hungary [13]. This paper summarized the abilities of the different planners' services, focusing on multimodality.

Moreover, not just the algorithm, but the hardware is also important. In the first period of modern computer technology, transistor miniaturization was the solution to increase the computing capability. This was the basis of Moore's law which was originally a prediction by Gordon Moore [14]. He was the co-founder of Fairchild Semiconductor Inc. and Intel corporations. In the last decades of the 20th century, computer technology started to reach an upper limit of the possible computing speed, what was that time the clock frequency of the transistors, expressed in Hz. Therefore, researchers were looking for another possible option. The solution was the ability of computing simultaneously, also called multi-threading [15].

The Ant Colony algorithm is widely used and researched in the Civil Engineering area. The algorithm was applied to structural calculation by Anikó Csébfalvi [16] [17] [18] [19] [20]. In these papers, the Ant Colony algorithm (ACO) was combined with the Genetic algorithm (GA) and the local search (LS) solution. The method called Angel was based on the evaluation of the approaches ACO, GA and LS. This complex approach resulted in significant time saving, and generated a more stable result. With this solution, the elastic-plastic material load-deflection can be handled. Moreover, a research was made to compare the capability of hybrid metaheuristic methods [21] and also hyperheuristic methods [22]. A special usage of these methods was a project scheduling case to support a water management program [23]. A deeper research of the Angel algorithm leads to a new approach by Anikó Csébfalvi [24] [25]. This possible combination was also researched by Ali Kaveh and Siamak Talatahari [26]. In their article, the heuristic particle swarm ant colony optimization (HPSACO) was introduced. Thanks to the development, the improved algorithm was more efficient and more robust; therefore, higher convergence can be reached than with former solutions. The aim of the research was to achieve a more effective tool to support the calculation for the optimal design of trusses. Another aim was the development of the design space's extensions. In this approach, a partitioning technique is required to avoid the fall of the searching speed [27].

A parallel computing technique was examined by Péter Iványi [28, 29] [30] [31] [32] [33] [34] [35] [36] [37] [38] [39] [40][41][42][43][44]. The aim was to use the modern multithreaded processors' capability to increase the computational speed. The research scope was parallel computing, later on cloud computing; and, last but not least, CUDA (Compute Unified Device Architecture) programming. CUDA is a parallel computing interface developed by Nvidia [45]. With this technique, the GPU's (Graphics Processing Unit) more effective parallel computing capability can be used instead, or, in this case, it can supplement the less effective CPU (Central Process Unit). This technique is also used in Discrete Element Method (DEM) [46]. The parallel and asynchronous grid computation methods are also researched in connection with DEM for American options derivatives [47] and for the simulation of 3D electrophoresis [48].

The routing problem usually needs to take multiple criteria into account. There are plenty of solutions conceivable in addition to the methods listed above. Heinz Spiess and Michael Florian [69] described a possible method to formulate the transit assignment problem to the linear programming problem. Moreover, non-linear cost extension is also taken into account.

The algorithm for real time network was developed by Jerald Jariyasunant, Eric Mai and Raja Sengupta [49]. Here, the main technique was the combination of a pre-calculated network and the real dataset. With the pre-calculation, sufficient computing time could be reached.

A routing program prototype was made by Man-chun Tan, C.O. Tong, S.C. Wong and Jian-min Xu in Hong Kong and Guangzhou area [50]. The main expectations were acceptable time and the transfer-walk criterion during the routing process.

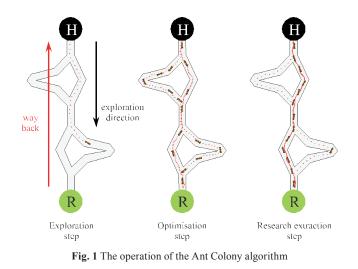
A modular transportation approach was presented by Péter Bocz, Lajos Kisgyörgy and Gergely Vasvári [51]. In this paper, a demand-controlled system with a statistical impact was described. Another approach is stochastic traffic assignment [52]. The imitation of the transport system was the aim with the Stochastic User Equilibrium (SUE) algorithm. With this technique, the condition and performance of the network can be rapidly determined. The logit model [53] was also considered.

In a wider scope, there are plenty other algorithms in the Civil Engineering area that can also make up a good basis for routing. A good approach could be Enhanced Colliding Bodies Optimization (ECBO), which is used for size optimization in the case of monopole steel structures [54]. Another research scope is the Volume-constrained Expected Compliance Minimization with Probabilistic Loading Directions supported by the optimal solution searching optimality criteria (OC) method [55]. An efficient approach is the shortest route trees, that helps the meshing procedure [56].

Calculation time is a key factor during the routing process, especially in a multimodal case, which results in a more complex network. Moreover, it is also worth mentioning Y. Sheffi's work [6] with the extension of the autonomous vehicles [57]. The smart city [58] conception should also be taken into account, which assumes a higher level of cooperation (IoT^1) or control.

One of the main focuses of autonomous car development is to implement more and more effective algorithms for routing. This is based on the fact, that the main computer needs to process plenty of sensor signals to create a special model for the driving algorithm. This process requires enormous computation capacity which causes high heat emission and requires large space. On one hand, this issue can be handled by the processor manufacturers [59] [60] [61] [62] [63]; however, an accurate and fast algorithm is also needed, which can use the parallel computing capability of these processors.

Based on our previous research and the studies mentioned above, the hypothesis is the following: if the iteration number of a heuristic route planning algorithm (in this case the Ant Colony) is decreased and the parallel computing option is used instead of the iteration, the total time usage decreases. In addition, the resource utilization increases. In this study, an ant corresponds to a processor's thread. The innovation of the paper is the parallel computational Ant Colony algorithm and the adequate parameter setting for parallel computation based on test results.



2 Methodology – Implementation of the Ant Colony algorithm

The Ant Colony algorithm was invented by Marco Dorigo [64] [65]. This method simulates an ant colony's resource gathering process. The aim of the hive (colony) is to find enough resource located in attainable distance. The ants developed a method to serve the effective operation of the hive. In every case, an explorer ant is sent. If the ant finds something needed by the hive, the route is signed with pheromone vapor on the way back. The next ant is able to follow this sign or choose another direction. The stronger the pheromone sign, the more ants follow. It is easily understandable that closer or bigger the resource is, the stronger the sign is. If the path is used less frequently, the pheromone sign evaporates. The following figure helps to understand this method (Fig. 1), where H is the hive and R is the resource.

One of our previous papers [1] investigated the related research works in connection with this field, and laid down the methodological basis of the current work. The model is based on choosing the probability of an actual connection [66]. The essence of the method is the pheromone change that indicates the level of suitability of the path.

$$\Delta \tau_{ijk} = \frac{Q}{d_k} \tag{1}$$

In Eq. (1), $\Delta \tau_{ijk}$ indicates the pheromone change, $[d_k] = (m)$ is a route length generated by the actual ants, and [Q] = (m) is a factor to control the importance of the distance. *i* and *j* are identifiers of the node. *k* is the identifier of the ant.

The next step is to collect the pheromone changes from the individual ants.

¹ Internet of Things

$$\Sigma \Delta \tau_{ij} = \Sigma \Delta \tau_{ij} + \Sigma \tau_{ijk} \tag{2}$$

 $\Sigma \Delta \tau_{ij}$ is the summarized pheromone change between *i* and *j*. Thus, the following equation describes the update process:

$$\tau_{ij} = \tau_{ij} \cdot (1 - \varrho) + \Sigma \Delta \tau_{ij} \tag{3}$$

This equation summarizes the ants' path information and the evaporation of the pheromone. Here, $[\varrho] = (-)$ is the evaporation factor and $[\tau_{ij}] = (-)$ is the actual quantity of the pheromone.

Based on the previous equation, the choosing probability can be calculated as

$$p_{ij} = \begin{cases} \frac{\left(\tau_{ij}\right)^{\alpha} \left(\eta_{ij}\right)^{\beta}}{\sum\limits_{\substack{u \notin M_k}} \left(\tau_{ij}\right)^{\alpha} \left(\eta_{ij}\right)^{\beta}} & \text{if } j \notin M_k \\ 0 & \text{if } j \in M_k \end{cases}$$
(4)

where $[\eta_{ij}] = (1/m)$ is the inverse of the measured length between *i* and *j*, $[\alpha] = (-)$ is the weight of the pheromone, $[\beta] = (-)$ is the weight of the distance and M_k is the route of the specified ant.

At the beginning, the network is unknown, and it is discovered by the ant movements. By increasing the efficiency of the convergence, an elite ant is chosen. This feature is shown in the following figure, where the structogram [67] representation technique was used (Fig. 2).

Some protections are implemented to the algorithm, which help to avoid backtracking (dead end) and loops.

To make sure that the algorithm finds a solution, another feature was built in. If an ant is stuck in a dead end, it steps back to the last junction. This procedure is shown below (Fig. 3, Fig. 4, Fig. 5)

A dead end black list was made, in which the nodes with no further connections can be stored. If a node is stored in this list, the possibility of choosing that is changed to zero. This way, a sufficient solution and acceptable calculation time can be reached. If the destination is unreachable, an infinite distance is given. In 1965, two scientists presented a similar algorithm [68]. The result of the first case (Fig. 3, Table 1), the second case (Fig. 4, Table 2) and the third case (Fig. 5, Table 3) can be found below.

$length_{min} > length_{actual}$				
	/			
$length_{min} = length_{actual}$				
$length_{min} = length_{actual}$ $rute_{best} = route_{actual}$ $best_{ant} = ant_{actual}$				
$best_{ant} = ant_{actual}$				



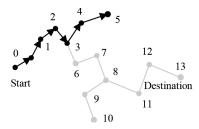


Fig. 3 Example network - First case: dead end was found

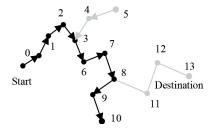


Fig. 4 Example network - Second case: dead end was found

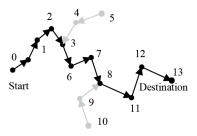


Fig. 5 Example network -Third case: destination was found

Table 1 The result of the first case

Route	0, 1, 2, 3, 4, 5					
Dead end list	-					
	Table 2 The result of the second case					
Route	0, 1, 2, 3, 6, 7, 8, 9, 10					
Dead end list	4, 5					
	Table 3 The result of the third case					
Route	0, 1, 2, 3, 6, 7, 8, 11, 12, 13					
Dead end list	4, 5, 9, 10					

The aim and also the novelty in this paper is speeding up the algorithm with parallel computing. This idea is based on nature and also on modern computers. There are several ants in nature which explore the environment around the hive. This can be simulated with multi-threaded systems. This article is in connection with a previously issued conference paper. [1].

The overview of the multi-threaded code is shown in Fig. 6, where each thread belongs to an Ant. The network is discovered by the ants, and the route attributes (Δt) are shared with each other after they return back to the hive.

Set th	he initial parameters: α , β , ϱ , Q ,	Nr0fAnt, iterations							
For <i>i</i>	terations = 0 to iteration models	ax do							
	Starting NrOfAnt pieces of processing threads								
	Route planning of Ant 1	Route planning of Ant 2		Route planning of Ant NrOfAnt					
	Calculate the pheromone matrix changes $\Delta \tau_1$ of Ant 1	Calculate the pheromone matrix changes $\Delta \tau_2$ of Ant 2		Calculate the pheromone matrix changes $\Delta \tau_{NrOfAnt}$ of Ant NrOfAnt					
	Update the pheromone matrix $\underline{\underline{\tau}} = \underline{\underline{\tau}} \cdot (1 - \varrho) + \Sigma \underline{\Delta \underline{\tau}}$								
	Generate the graphics of the $\underline{\tau}$								

Fig. 6 Structogram of the parallel Ant Colony algorithm

The pheromone matrix is updated after each iteration. The result of this method is a decreased number of iterations and time saving with the parallel computation, where each thread belongs to an Ant.

3 Results

3.1 Application on test network

The database implements the structure of the GTFS [General Transit Feed Specification]. There are public transport routes from Budapest and from Vienna. Some of the possible connections are stored, e.g.: railroads, bus and air lines. These are listed in Table 4 and shown on a network map (Fig. 7).

The GTFS data contains valid data from different transport operators and the links are shown in Fig. 7.

3.2 Initialization of the test

As outlined by the data, the algorithm is influenced by several factors that can modify the result. Therefore, a parameter study was implemented to examine the connection between the results and the factors. A HP ProLaint ML310e Gen8 v2 was used for the tests. The computer has an Intel Xeon E3-1220 CPU which runs on 3.10 GHz, and has 4 GB of system memory. The application was written in Java which runs on Ubuntu 14.04.4 LTS with OpenJDK Runtime Environment version 1.7.0_95. The following parameters were reviewed: α , β , ϱ , Q, NrOfAnt, Threads, iterations (Table 5). Parameters α , β , ϱ , Q modify the operation of the algorithm, and the other two affect the efficiency.

Furthermore, the execution time of the algorithm was also registered. All ants were supervised in every execution and in each iteration from the point it has been created

Destination Route						
M1	Mexikói út M / Vörösmarty tér					
M2	Örs vezér tere M+H / Déli pályaudvar M					
M3	Kőbánya-Kispest M / Újpest-központ M					
4	ÚjBuda-központ / Széll Kálmán tér M					
6	Móricz Zsigmond körtér / Széll Kálmán tér M					
7	Bosnyák tér / Albertfalva vasútállomás					
173	Bornemissza tér vá. / Újpalota, Nyírpalota út					
200E	Határ út M / Liszt Ferenc Airport 2					
Railjet	Budapest-Keleti / Wien Westbahnhof					
912/942	Budapest-Keleti / Sopron					
2821	Sopron / Wien Meidling					
130	Budapest-Keleti / Bratislava Hlavná Stanica					
2513	Bratislava Hlavná Stanica / Wien Hauptbahnhof					
U1	Reumann Platz / Leopoldau					
U2	Karlsplatz / Aspernstraße					
U3	Ottakring / Simmering					
U4	Hütteldorf / Heiligenstadt					
U6	Siebenhirten / Floridsdorf					
CAT	Bahnhof Wien Mitte / Flughafen Wien					
S7	Rennweg / Flughafen Wien					
Volán	Budapest, Népliget aut. pu. / Wien, VIB, U3 Station Erdberg					
OS 714	BUD Nemzetközi repülőtér / Flughafen Wien					

Table 4 Transport trips in GTES

until it reached the defined target. Additionally, the time of the iteration was also registered. This data will be important in the future, when parallel operation will be reviewed. A TOP 5 list was created from the results of the shortest route length. The number of occurrences was counted to each of the lengths in every case. The parameters were evaluated on the basis of the number of occurrences. First, parameters *NrOfAnt* and *Threads* were reviewed (Fig. 8).

Table 5 Values of parameter study						
Parameter	Value					
α	Pheromone importance indicator	0.1; 0.6; 1.1; 5.1				
β	Distance importance indicator	0.05; 0.15; 0.45				
6	Control the loss of pheromone	0.1; 0.2; 0.3; 0.9				
Q	Control the distance	100; 600; 5100				
NrOfAnt	Number of ants	1; 2; 8				
Threads	Number of threads	1; 2; 8				
iterations	Number of iterations	10				

The points (Fig. 8) were the result of the number of ants multiplied by the measured distance. On the X-axis, the number of ants and the number of threads used are shown, and the Y-axis represents the calculated points. Where the number of threads is smaller than the number of ants, the ants wait until a thread becomes free. Every *NrOfAnt* and *Threads* settings are executed with the same α , β , ρ , Q parameters. After a closer look at the results, it can be declared that using only 1 thread results in the highest points until 6 ants. Consequently, the higher number of threads has an effect over 7 ants. Using these settings for the algorithm, the result is stably TOP 1 in approx. 40% of the cases from the sum of the TOP 5 results, which is a good outcome.

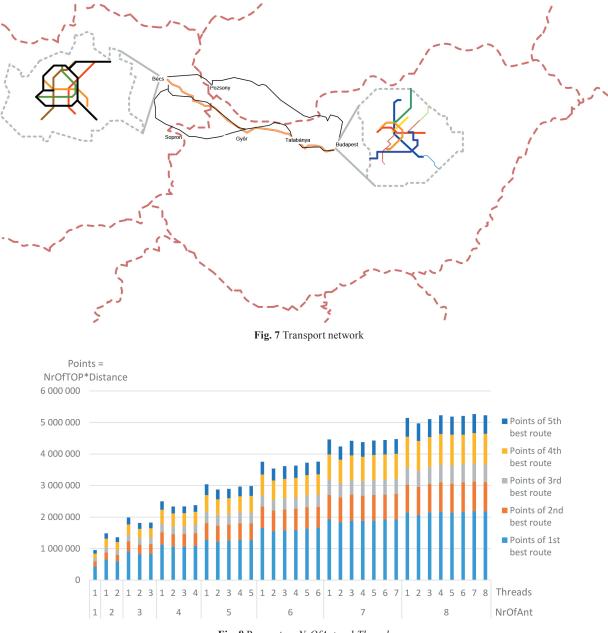


Fig. 8 Parameters NrOfAnt and Threads

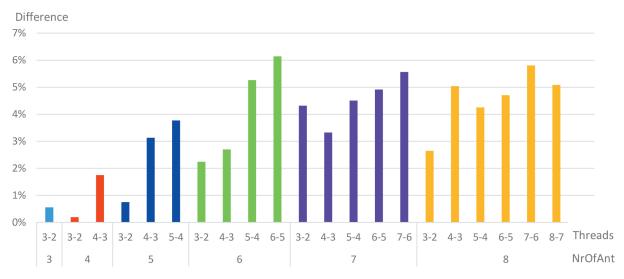
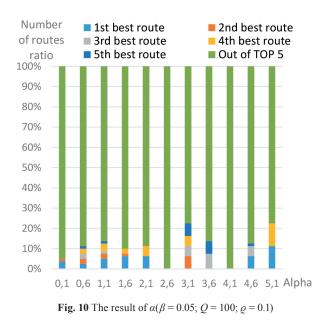
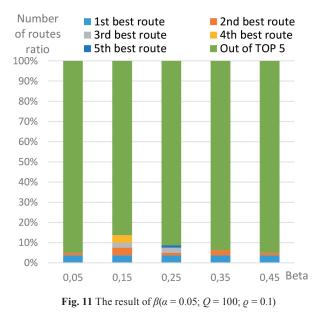


Fig. 9 Cumulated Ratio of point differences





Continuing the data analysis, the steps are also important between the different settings. The results are presented in Fig. 9.

The diagram indicates the cumulated ratio differences in a similar *NrOfAnt* group. Based on this, it is worth raising the number of threads, because the cumulated thread differences are always positive. And at *NrOfAnt* = 8, the result is not just relatively positive, but it is absolutely positive, too (Fig. 8).

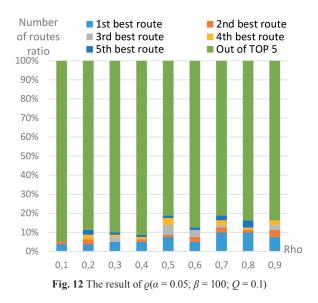
Consequently, in the next step with *NrOfAnt* = 8 and *Threads* = 8 settings, the α , β , ρ , Q parameters are investigated. In the following figure (Fig. 10), the α parameter results are presented. The horizontal axis shows the alpha value and the vertical axis represents the number of the

routes ratio. So for example the $\alpha = 0.1$ settings eventuate approx. 4% TOP 1 and 1% TOP 2 routes. So, based on this, parameters $\alpha = 3.1$ and $\alpha = 5.1$ cause the highest number of TOP routes.

Fig. 11 presents the β parameter results. Based on my previous investigation, β is set between 0.05 and 0.45, because it has not been effective over 0.5. [1] Based on this result, β under 0.25 is effective.

The result of ρ is presented in Fig. 12. The highest ratio of the best route can be found between 0.5 and 0.9. Raising the value of ρ is causing higher number of the TOP 5 result.

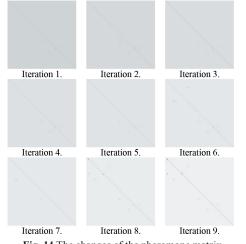
Last, but not least, results for Q are in the figure above (Fig. 13). The diagram shows that Q over 1600 causes more TOP results.

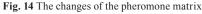


4 Analysis and discussion

Each result is calculated after 10 iterations. The pheromone matrix ($\underline{\tau}$) was updated after every step. This can be graphically seen in the following pictures (Fig. 14). One can see 9 examples from the 10 iterations below.

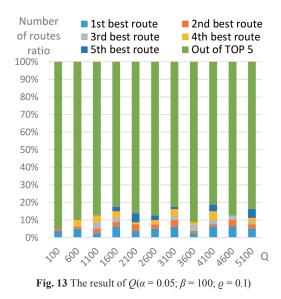
In the figure, darker pixels mean stronger pheromone evaporation rates. The vertical axis represents the source and the horizontal axis represents the destination. The structure explanation of the matrix can be found in Fig. 15.





	1. column		<i>j</i> . column		n. column
1. row					
:					
<i>i</i> . row			<i>i., j.</i> pixel		
:					
n. row					
	-			-	nicture

Fig. 15 The graphics of the pheromone matrix structure



After the parameters were investigated, the overall outcome was tested. The goal is to find alternative routes along the best route. The result was summarized and filtered. Those parameters were collected by which all TOP 5 routes were found. Efficiency was calculated between all possibilities and the TOP 5 routes. A top list was made from the efficiency (Table 6), this list contains values over 20%.

Examining the data (Table 6) shows that only the value of β is fixed, the other values are less stable.

In the next step, the execution time was investigated (Fig. 16)

In the diagram, the average individual ant time and the average sum time is demonstrated. In the second y axis the ratio of the average sum and the average individual time was presented. The decrease in the ratio means that the sum and the individual average time are getting closer. Consequently, the meaning of the ratio is that the time gain is getting smaller with a higher number of ants.

5 Conclusions

In this study, the parallelization of the Ant Colony algorithm was introduced based on previous research results, for example in Civil Engineering area. Different parameters and thread counts were investigated. In conclusion, we can emphasize that it is worth using more threads, because the probability of finding valid routes is increasing. In addition, time usage and resource utilization are also better. This capability is beneficial in a large-sized network. Based on the results, the hypothesis is considered to be supported.





 Table 6 The best parameters of the Ant Colony algorithm

 (Iteration = 10, NrOfAnt = 8 and Threads = 8)

				-					s.	es	
ø	β	Ø	6	1st best route	2 nd best route	3rd best route	4 th best route	5 th best route	SUM of TOP :	Sum of all routes	Efficiency
3.1	0.05	0.3	600	9	5	3	7	1	25	80	31%
0.1	0.05	0.4	4100	9	4	2	1	3	19	80	24%
1.1	0.05	0.8	1600	9	3	2	1	1	16	80	20%
1.1	0.05	0.4	4100	8	3	2	2	1	16	80	20%
0.6	0.05	0.5	100	8	1	1	2	4	16	80	20%
4.1	0.05	0.2	1600	7	3	2	2	3	17	80	21%
1.1	0.05	0.7	1100	7	3	1	3	3	17	80	21%
5.1	0.05	0.6	1100	7	1	4	5	1	18	80	23%
1.1	0.05	0.5	2100	6	4	2	2	2	16	80	20%
0.6	0.05	0.6	5100	6	3	4	2	4	19	80	24%
1.1	0.05	0.9	2600	6	3	3	5	3	20	80	25%
0.6	0.05	0.2	2600	6	3	3	4	1	17	80	21%
1.6	0.05	0.8	1600	6	1	4	3	3	17	80	21%
1.1	0.05	0.6	2600	6	1	4	2	3	16	80	20%
2.1	0.05	0.6	5100	5	4	2	5	1	17	80	21%
0.6	0.05	0.4	100	5	4	2	2	5	18	80	23%
1.1	0.05	0.4	4600	5	3	1	4	4	17	80	21%
2.6	0.05	0.1	4100	5	2	4	3	2	16	80	20%
0.1	0.05	0.7	3600	5	1	3	3	5	17	80	21%
0.6	0.05	0.9	3100	4	6	3	3	1	17	80	21%
1.6	0.15	0.2	600	4	4	4	5	2	19	80	24%
5.1	0.05	0.7	2600	4	4	3	4	3	18	80	23%
0.6	0.05	0.6	3600	4	4	2	3	4	17	80	21%
0.6	0.05	0.2	2100	4	3	4	3	2	16	80	20%
1.6	0.05	0.4	5100	3	4	5	3	2	17	80	21%
2.1	0.05	0.8	3600	2	3	4	4	5	18	80	23%

In further research, this algorithm will be applied for large-sized networks on a multi-core system, where more threads will be available. Further development will be the determination of the full resilience function, which contains not just the distance used in this paper, but the travelling time and some other factors, too. Another research aim is to integrate autonomous vehicles in the routing process. This can be one supporting element for the Smart City conception and a central traffic management system.

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